

Architectural design: the coordination perspective

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Outline

1. Visual semantics for Reo

▶ Connector colouring (CC)¹

2. Locality (concurrency)

▶ partial connector colouring (PCC)²

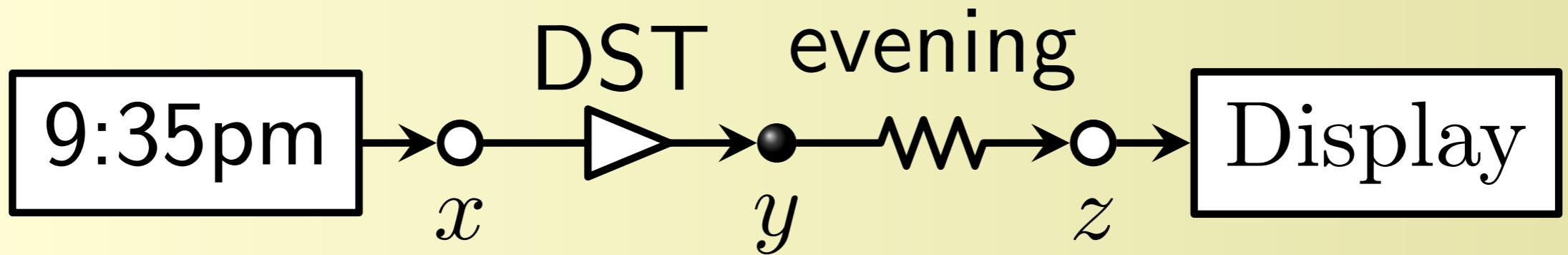
3. Constraints

▶ SAT solving with data for Reo³

¹ Dave Clarke, David Costa, and Farhad Arbab. Connector colouring I: Synchronisation and context dependency

² Dave Clarke and José Proença. Partial connector colouring

³ Dave Clarke, José Proença, Alexander Lazovik, and Farhad Arbab, Channel-based coordination via constraint satisfaction
José Proença, Dave Clarke, Interactive interaction constraints



$$x \rightarrow \hat{x} := 9:35\text{pm}$$

$$x \leftrightarrow y$$

$$y \rightarrow \hat{y} := \text{DST}(\hat{x})$$

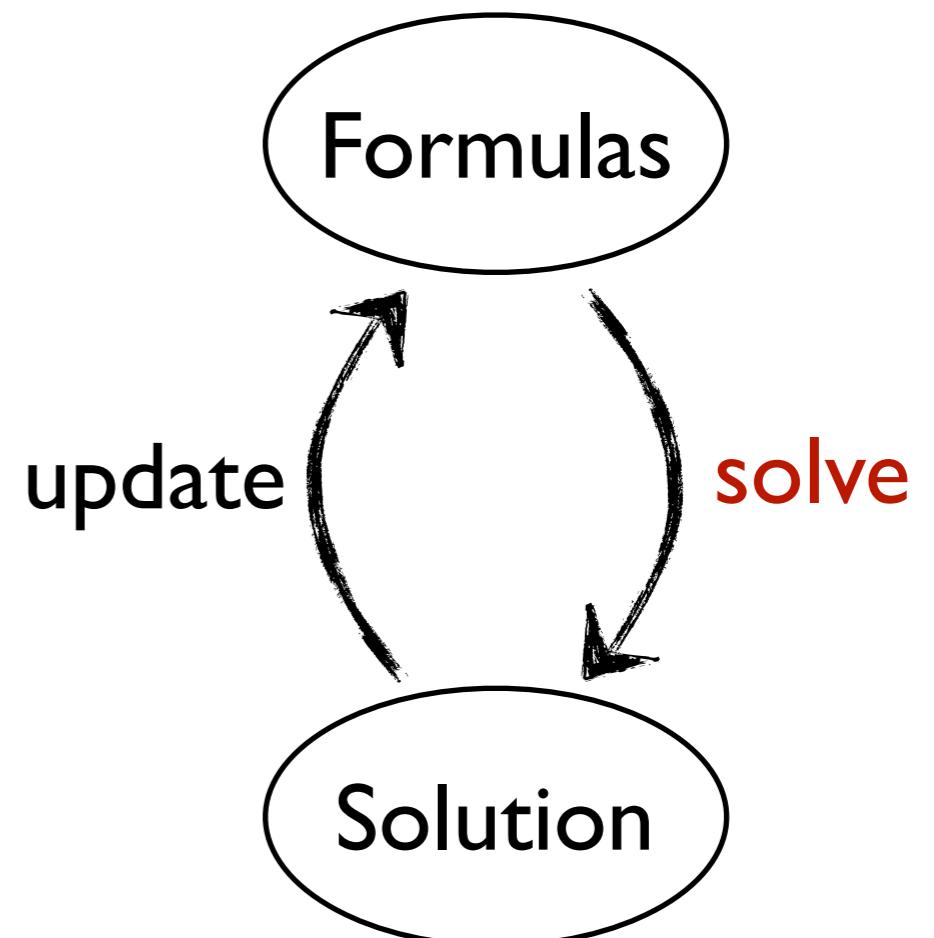
$$(y \wedge \text{evening}(\hat{y})) \leftrightarrow z$$

$$z \rightarrow \hat{z} := \hat{y}$$

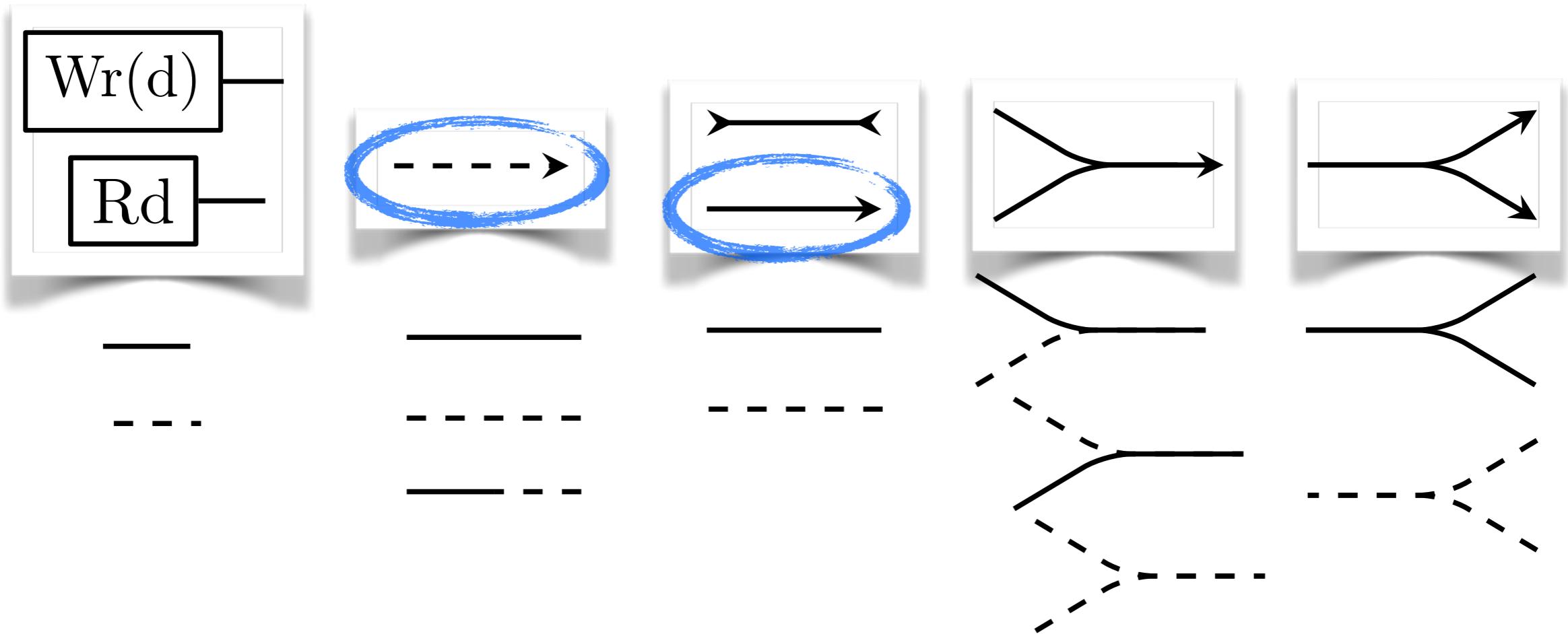
Reo as (Interactive) Constraints

Coordination as constraint satisfaction

$a \longrightarrow b$	$a \leftrightarrow b$ $b \rightarrow \hat{b} := \hat{a}$
$a \dashrightarrow b$	$b \rightarrow a$ $b \rightarrow \hat{b} := \hat{a}$
$a \xrightarrow{P} b$	$b \rightarrow \hat{b} := \hat{a}$ $(a \wedge P(\hat{a})) \leftrightarrow b$
$a \xrightarrow{f} b$	$a \leftrightarrow b$ $b \rightarrow \hat{b} := f(\hat{a})$



Exercise: write constraints



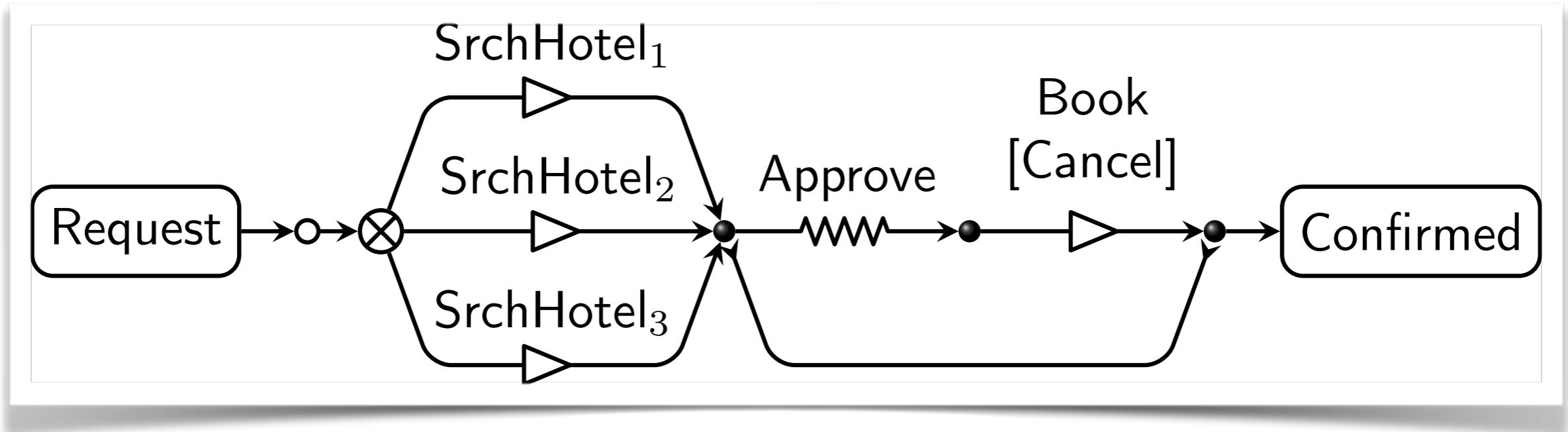
Building up constraints

- Connector colouring as constraints
- Data constraints
- Interactive constraints

Context

(need for extra interaction)

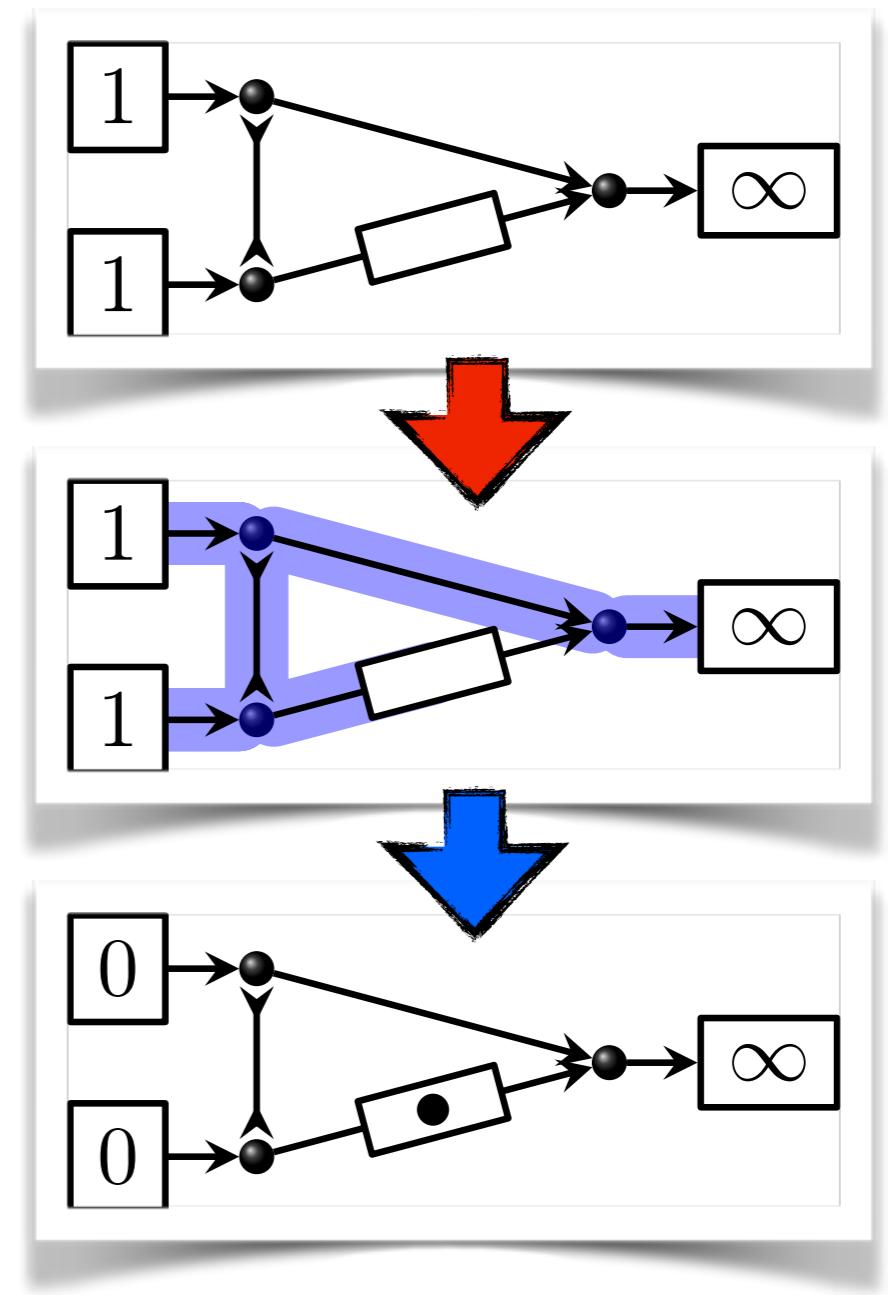
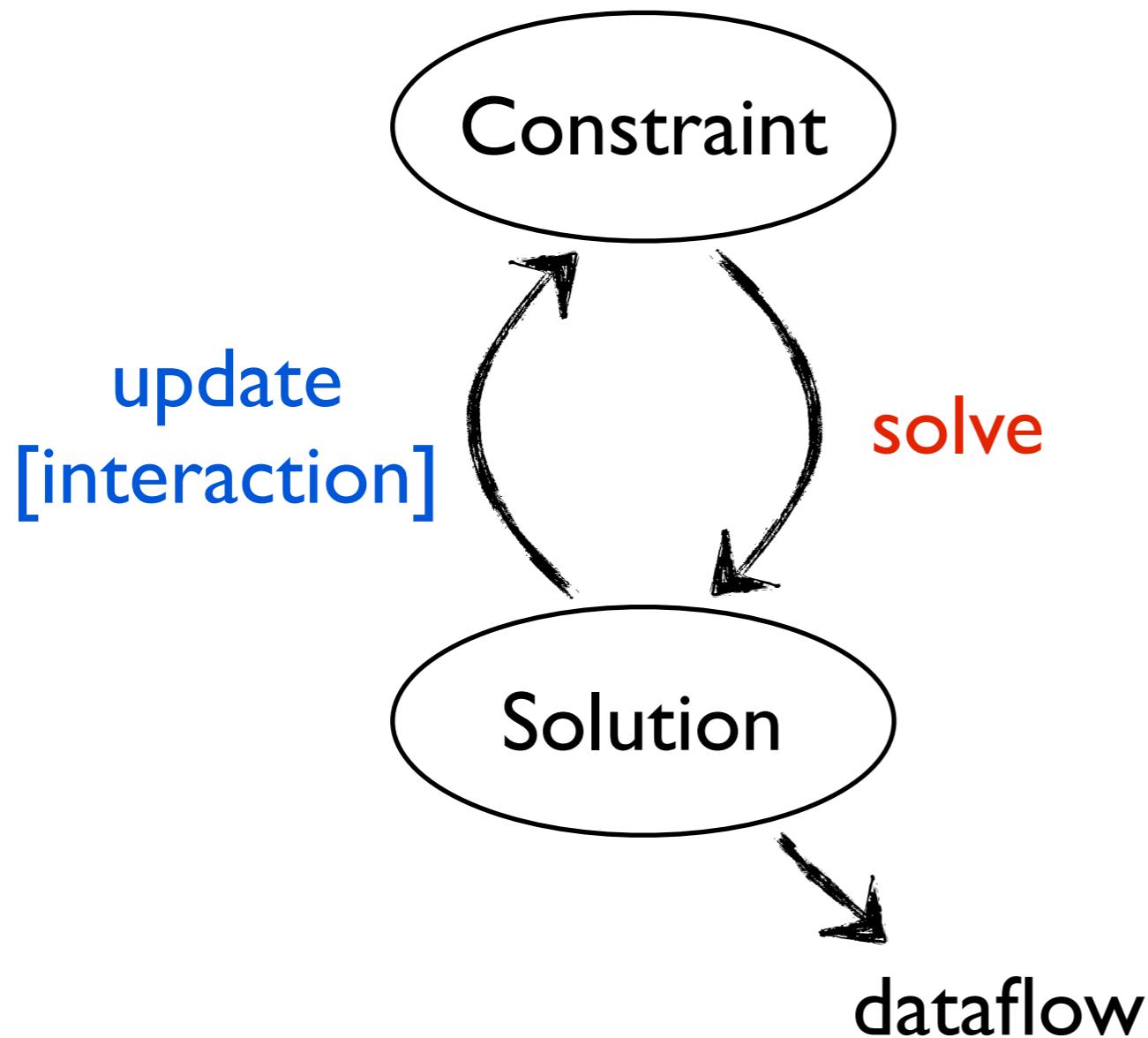
Hotel booking



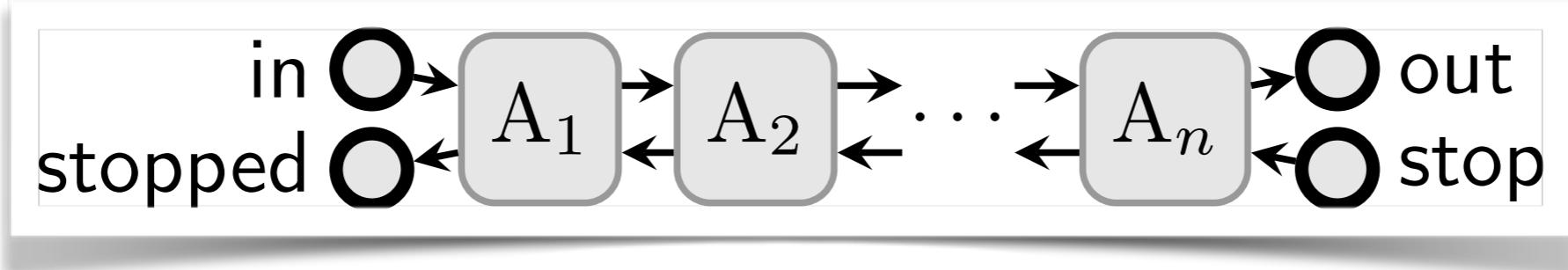
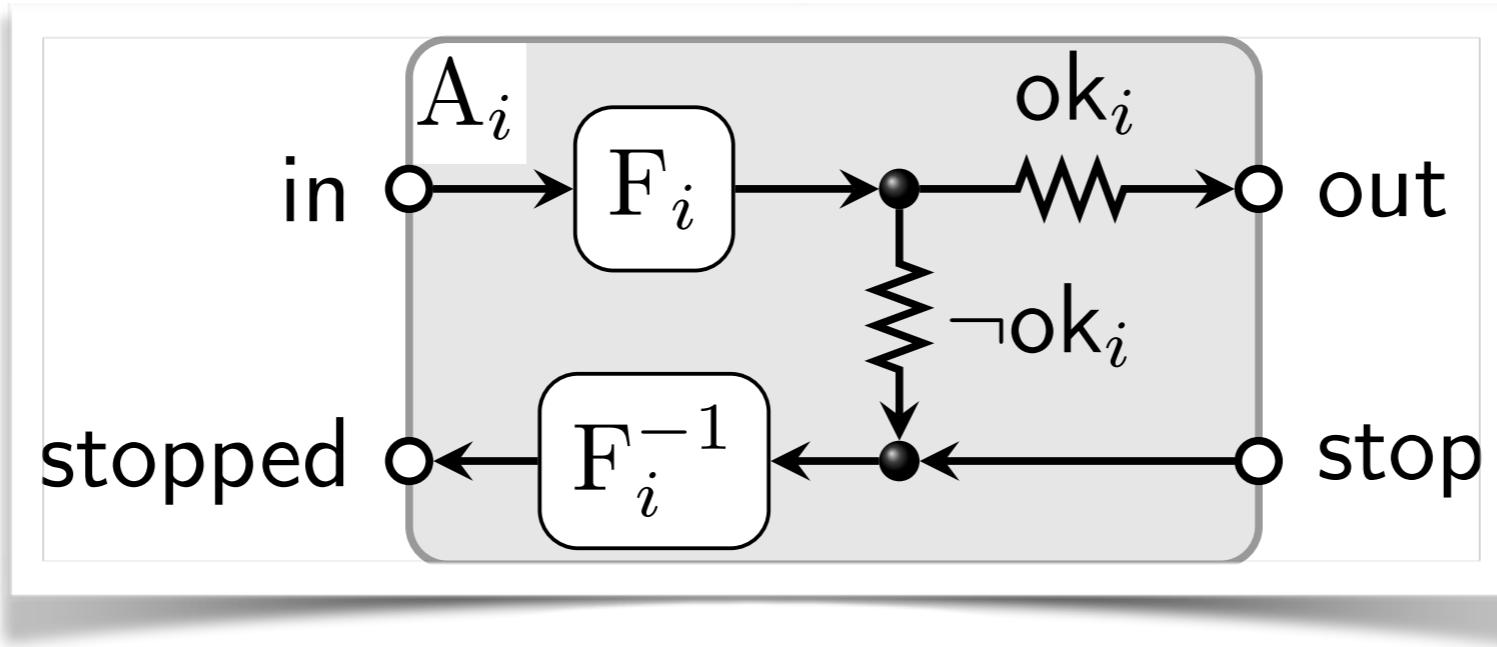
- Interaction with hotel repositories
- Interaction with users
- Interaction with hotels (availability & payment)

Problem of Interaction

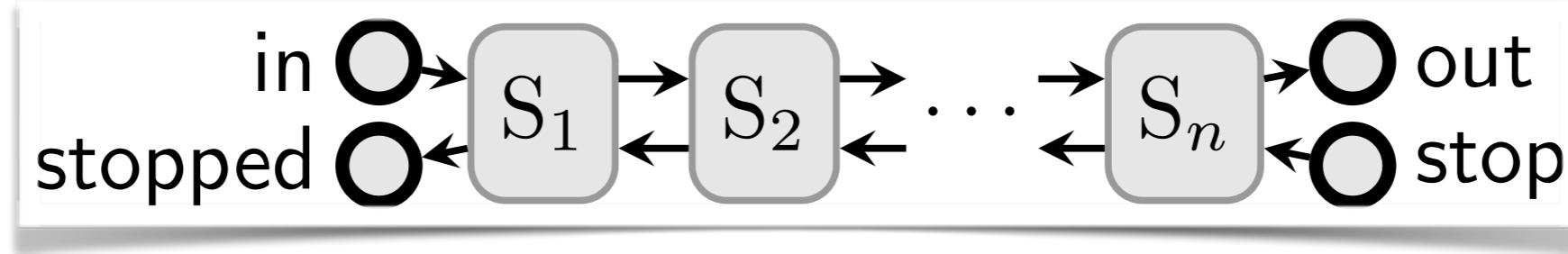
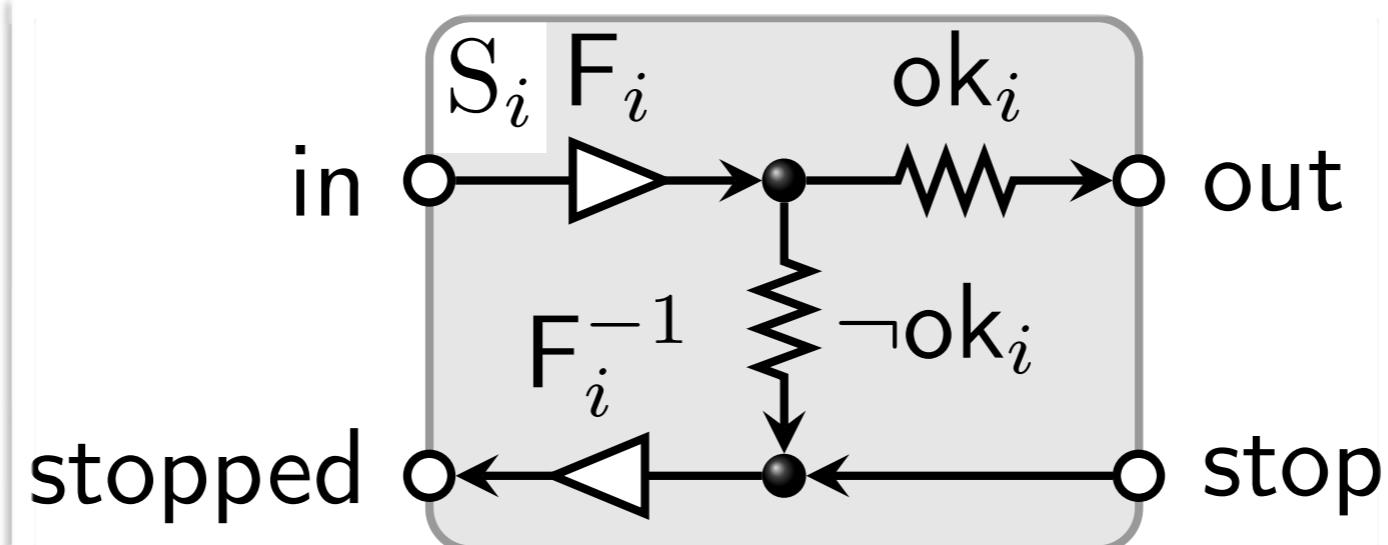
Coordination as constraint satisfaction



Asynchronous transactions



Synchronous transactions



Problem

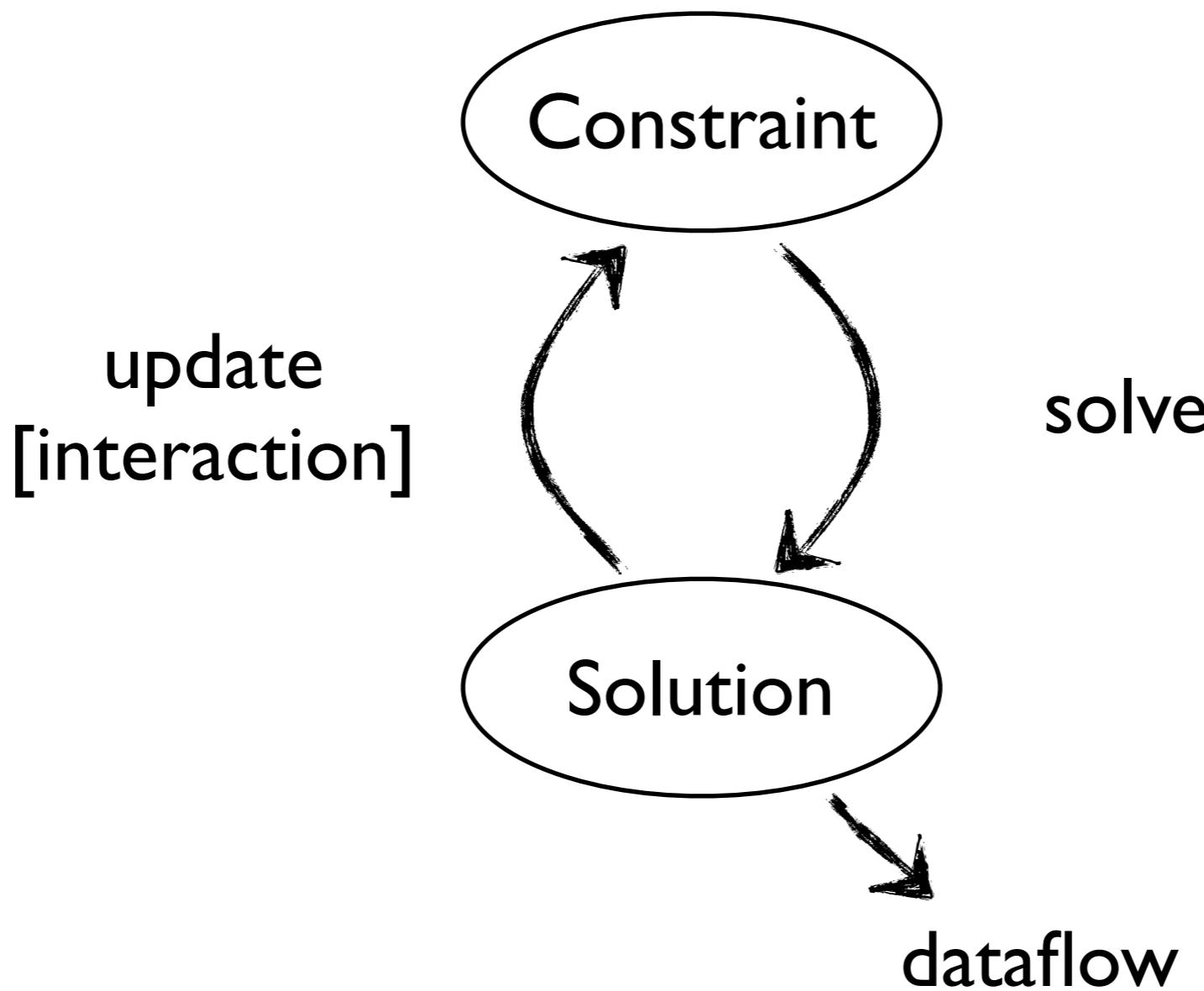
Interaction or Synchrony

but not both

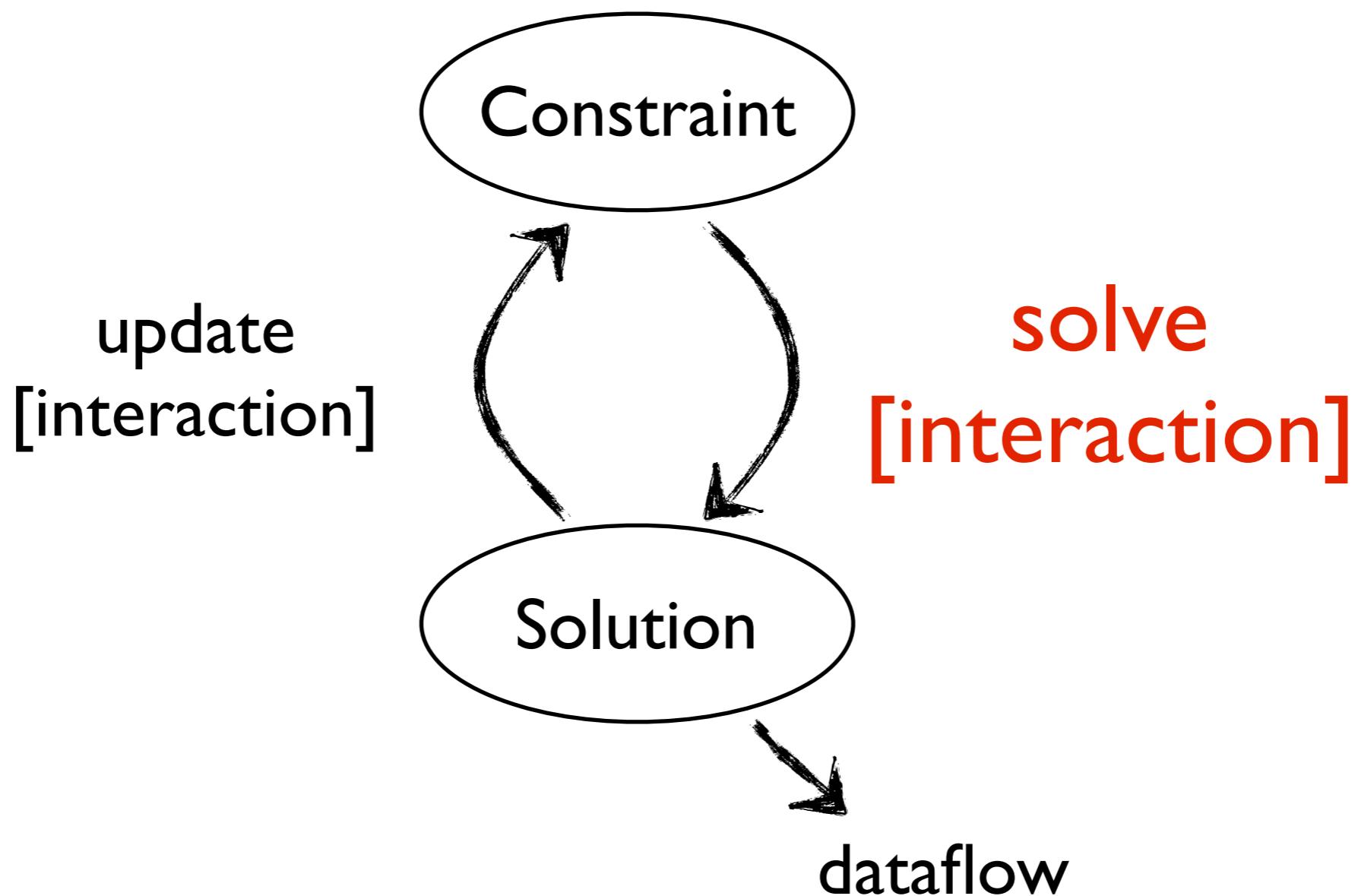
Solution

Interactive Interaction Constraints

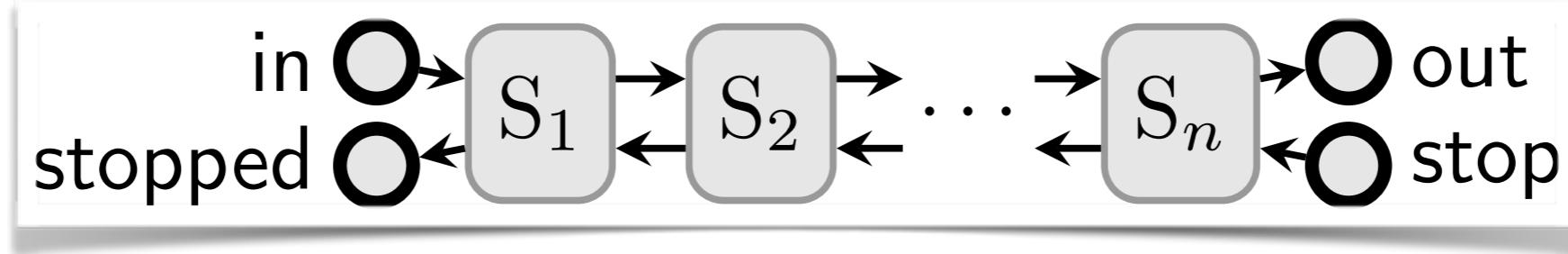
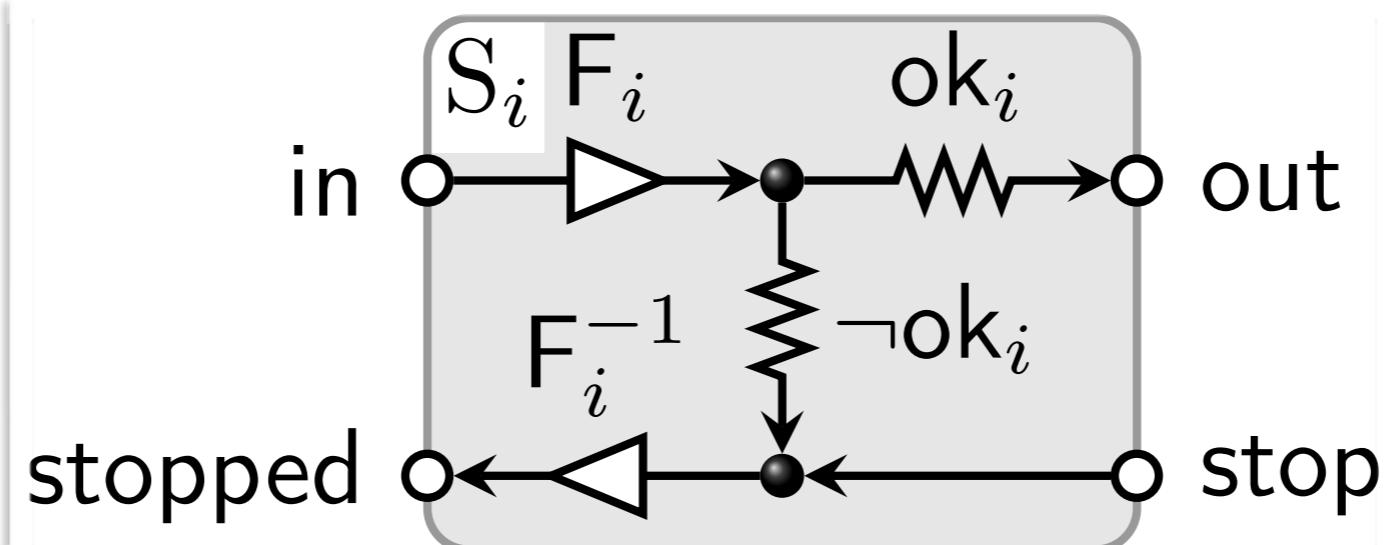
Coordination as constraint satisfaction



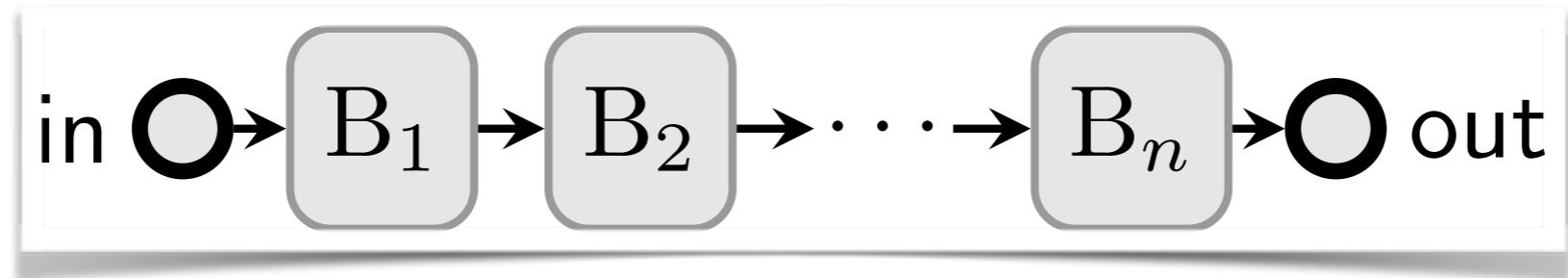
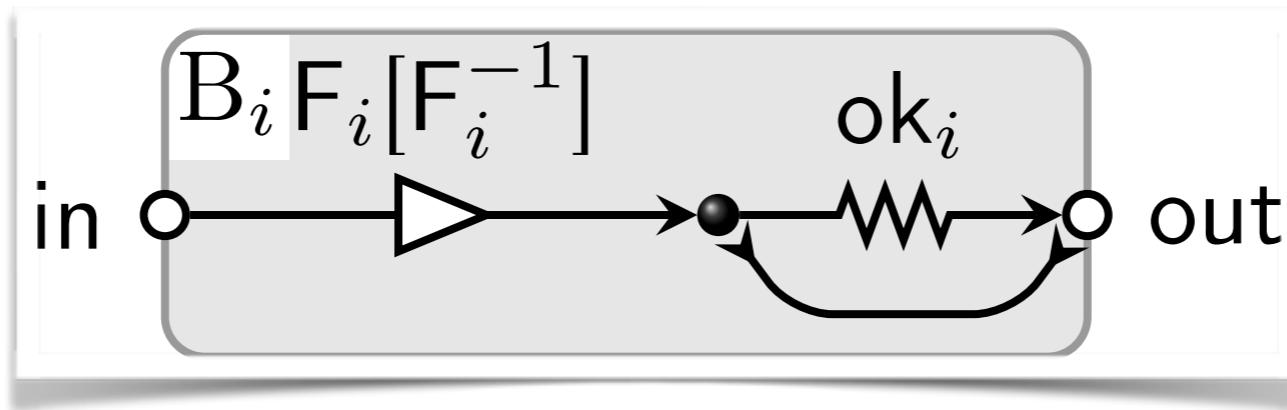
Coordination as **interactive** constraint satisfaction



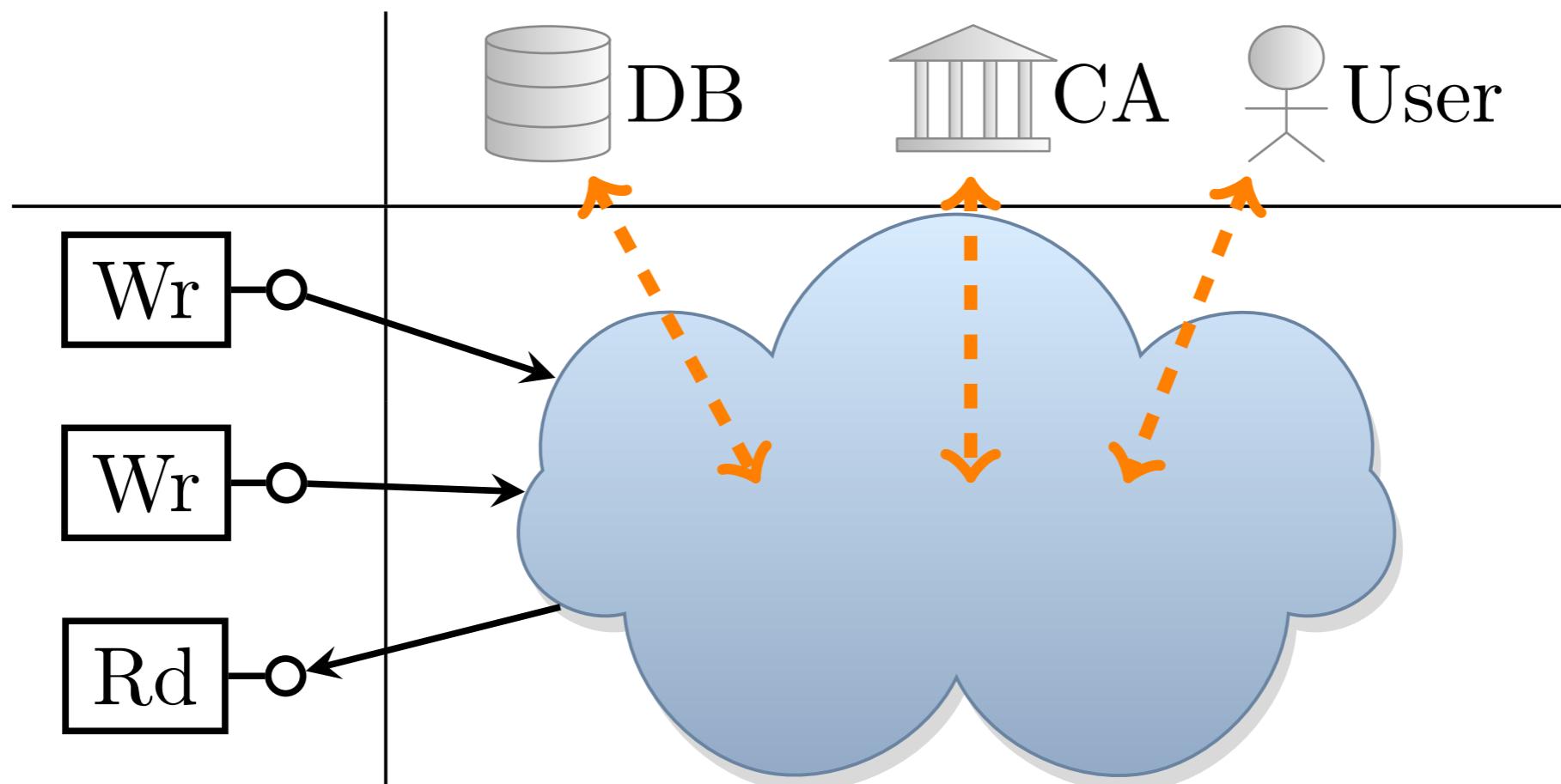
Synchronous transactions



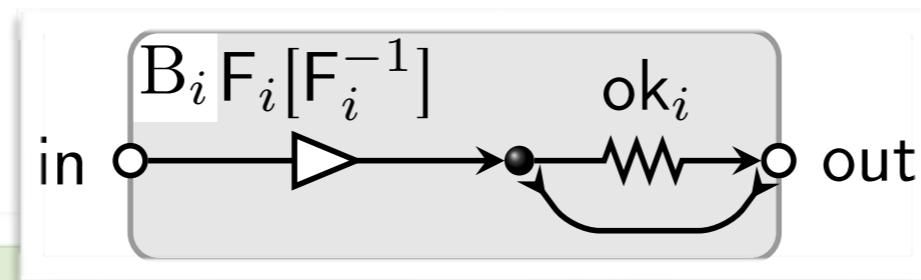
Synchronous transactions



Interactive Interaction Constraints

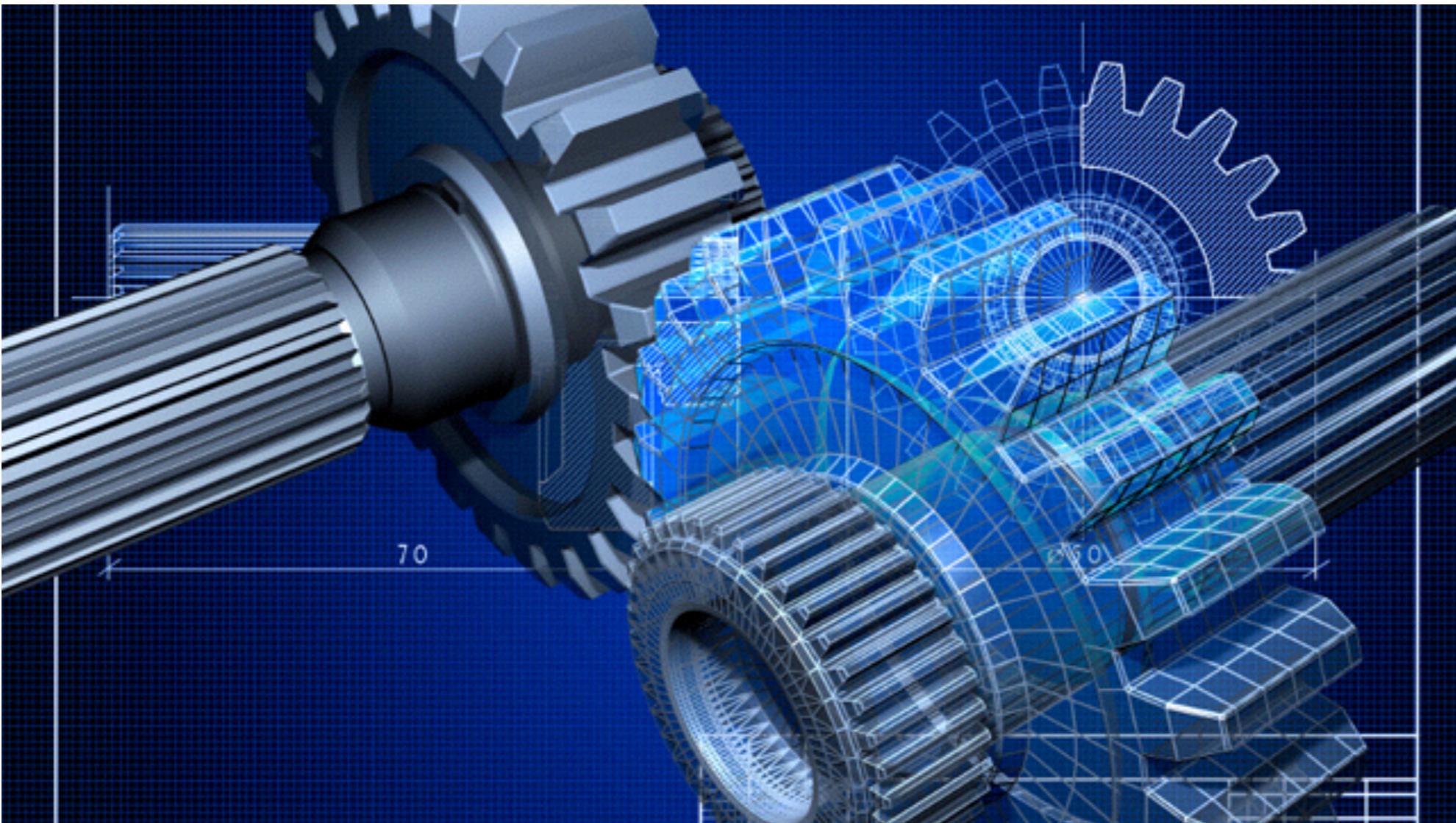


Scala / Java implementation



```
val f = Function("f") {  
    case s: String => /* do something */  
}  
  
val finv = Function("f^-1") {  
    case s: String => /* do something */  
}  
  
val ok = Predicate("ok") {  
    case s: String => /* do something */  
}  
  
val connector =  
    writer("in",List("a","b")) ++  
    transf("in","x",f,finv) ++  
    filter("x","out",ok) ++  
    sdrain("x","out") ++  
    reader("out",2)  
  
connector.run()
```

```
class Filter(a: String, b: String,  
            p: Predicate)  
extends ... {  
  
    def getConstraint = Constraint(  
        b --> a,  
        b --> (b := a),  
        b --> (a :< p),  
        (a /\ (a :< p)) --> b  
    )  
  
    /* override def update(s:  
        Option[Solution]) = ... */  
}
```



More about constraints

CC2 as constraints

- [*Colouring*: $\text{End} \rightarrow \{\text{Flow}, \text{NoFlow}\}$]
- *Formula*: Boolean over End
- *Composition* = conjunction

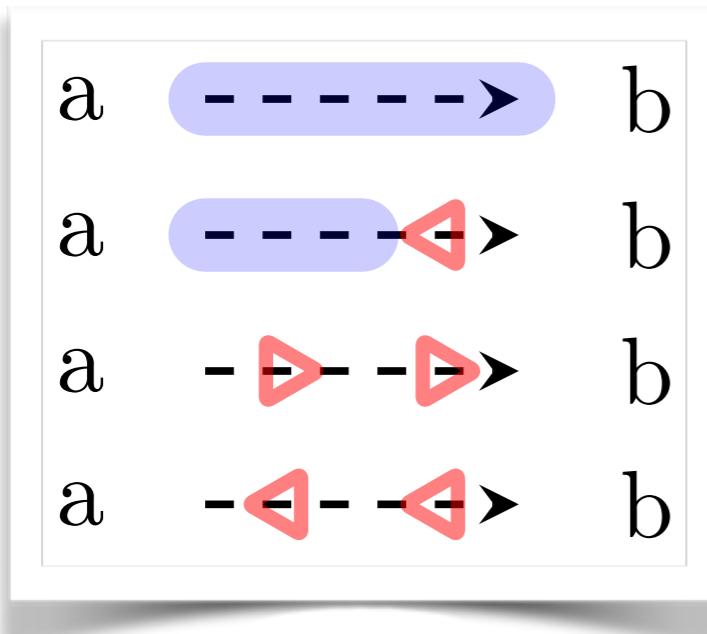
a \dashrightarrow b
a \dashrightarrow b
a \dashrightarrow b

$(\neg a \wedge \neg b) \vee$
 $(a \wedge b) \vee$
 $(a \wedge \neg b)$

$b \rightarrow a$

CC3 as constraints

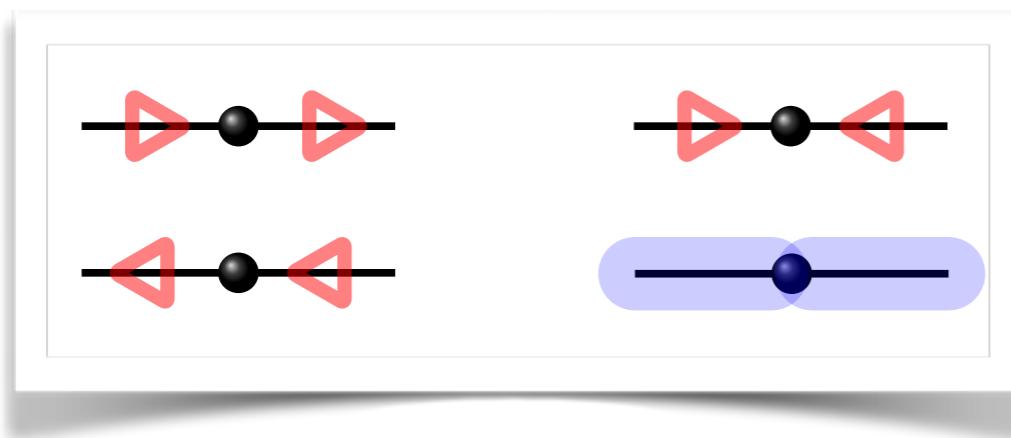
- [*Colouring*: $\text{End} \rightarrow \{\text{Flow}, \text{GiveReason}, \text{GetReason}\}$]
- *Formula*: Boolean over End , End_{src} , End_{snk}
- $a = \text{flow on } a$; $a_{\text{src}} = \text{give reason}$; $\neg b_{\text{snk}} = \text{get reason}$


$$\begin{aligned} & b \rightarrow a \wedge \\ & (a \wedge \neg b) \rightarrow (a \wedge \neg b_{\text{snk}}) \wedge \\ & \neg a \rightarrow (\neg b \wedge \neg a_{\text{src}} \wedge b_{\text{snk}}) \end{aligned}$$

CC3 as constraints

Composition

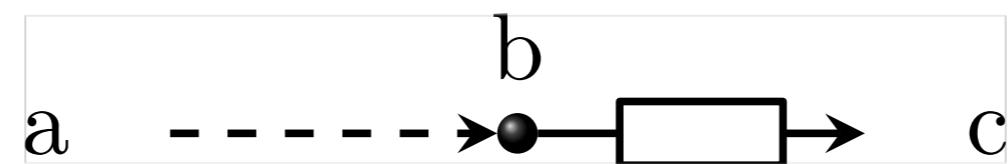
- a = flow on a ; a_{src} = give reason ; $\neg b_{snk}$ = get reason
- one-to-one composition: source to sink ends



$$\forall x \cdot x_{snk} \vee x_{src}$$

CC3 as constraints

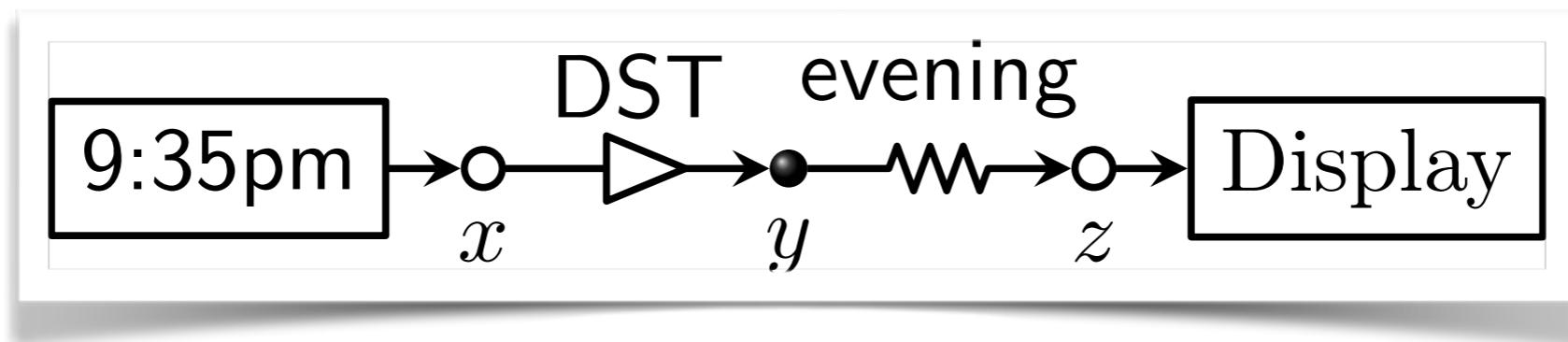
example



$$\begin{aligned}\phi = & \quad b \rightarrow a \wedge \neg c \\ & (a \wedge \neg b) \rightarrow (a \wedge \neg b_{snk}) \wedge \\ & \neg a \rightarrow (\neg b \wedge \neg a_{src} \wedge b_{snk}) \wedge \\ & (\neg b \rightarrow \neg b_{src}) \wedge c_{snk} \wedge \\ & b_{src} \vee b_{snk}\end{aligned}$$

$$\begin{aligned}\{a \wedge b \wedge \neg c \wedge c_{snk}\} &\models \phi \\ \{\neg a \wedge \neg b \wedge \neg c \wedge \neg a_{src} \wedge b_{snk} \wedge \neg b_{src} \wedge c_{snk}\} &\models \phi\end{aligned}$$

Data constraints



$$x \rightarrow \hat{x} := 9:35\text{pm}$$

$$x \leftrightarrow y$$

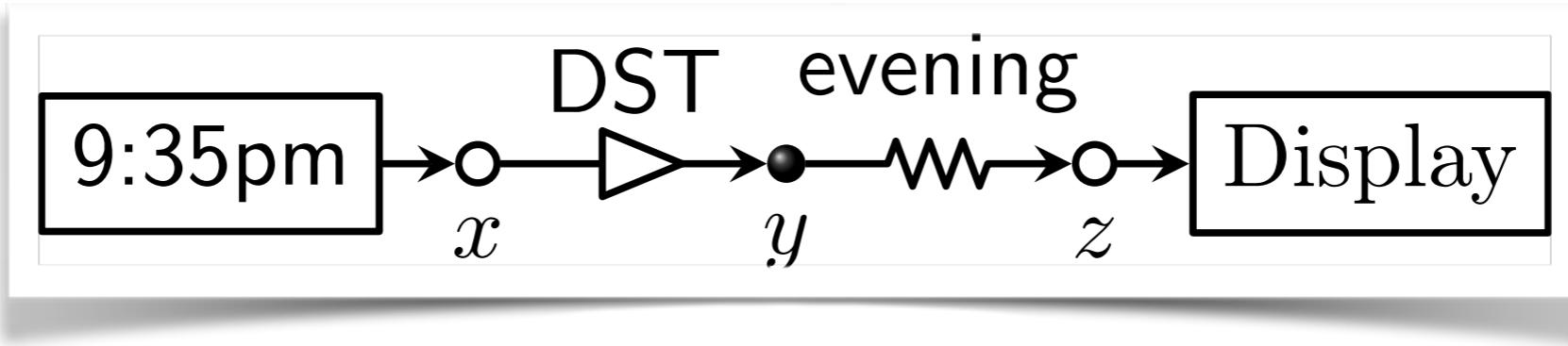
$$y \rightarrow \hat{y} := \text{DST}(\hat{x})$$

$$(y \wedge \text{evening}(\hat{y})) \leftrightarrow z$$

$$z \rightarrow \hat{z} := \hat{y}$$

How to solve this?

Predicate abstraction



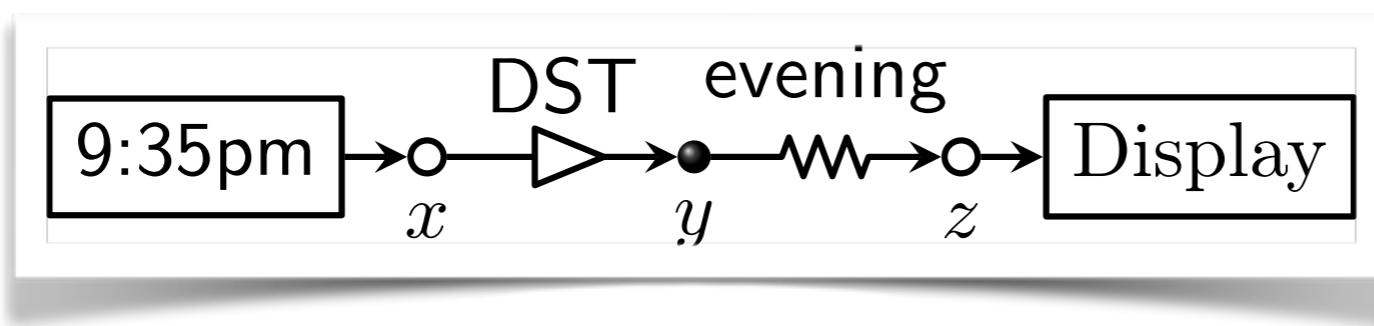
original

$$\begin{array}{lll} x \rightarrow \hat{x} := 9:35pm & x \leftrightarrow y & y \rightarrow \hat{y} := \text{DST}(\hat{x}) \\ (y \wedge \text{evening}(\hat{y})) \leftrightarrow z & & z \rightarrow \hat{z} := \hat{y} \end{array}$$

boolean

$$\begin{array}{ll} x \rightarrow \hat{x}_{\text{ev.dst}} := [\text{evening}(\text{DST}(9:35pm))] & x \leftrightarrow y \\ y \rightarrow \hat{y}_{\text{ev}} := \hat{x}_{\text{ev.dst}} & (y \wedge \hat{y}_{\text{ev}}) \leftrightarrow z \end{array}$$

Interaction via Choco



boolean

$$x \rightarrow \hat{x}_{\text{ev.dst}} := [\text{evening}(\text{DST}(9:35pm))] \quad x \leftrightarrow y$$
$$y \rightarrow \hat{y}_{\text{ev}} := \hat{x}_{\text{ev.dst}} \quad (y \wedge \hat{y}_{\text{ev}}) \leftrightarrow z$$

interactive

$$x \rightarrow \text{XPred}(\text{ev.dst}, x, 9:35pm) \quad x \leftrightarrow y$$
$$y \rightarrow \hat{y}_{\text{ev}} := \hat{x}_{\text{ev.dst}} \quad (y \wedge \hat{y}_{\text{ev}}) \leftrightarrow z$$

Interaction via Choco

- instance of a Choco constraint
- reacts when x or $\hat{x}_{\text{ev.dst}}$ is instantiated
- $\neg x \Rightarrow \hat{x}_{\text{ev.dst}}$ can be anything
- $x \Rightarrow \hat{x}_{\text{ev.dst}} = \text{ev}(\text{dst}(9:35\text{pm}))$

9:35pm

boolean

$x \rightarrow \hat{x}_{\text{ev.dst}}$

= [evening(DST(9:35pm))]

$x \leftrightarrow y$

$y \rightarrow \hat{y}_{\text{ev}} := \hat{x}_{\text{ev.dst}}$

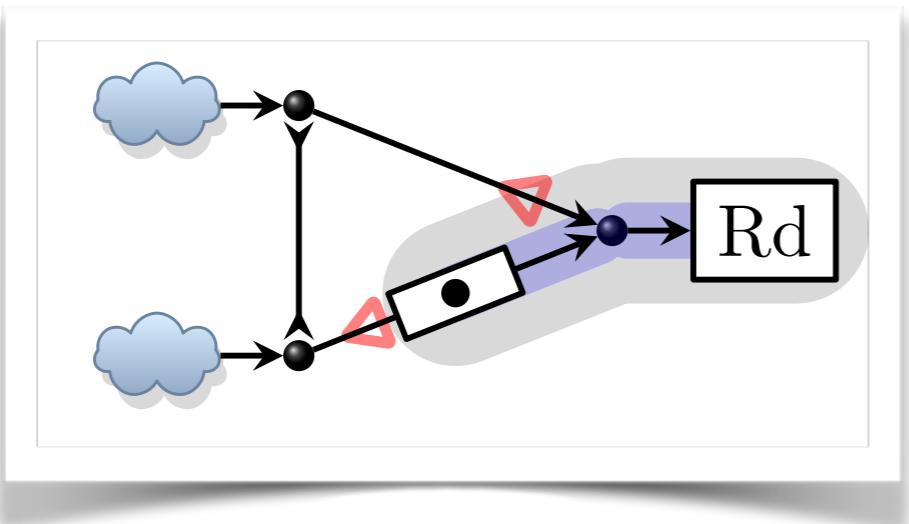
$(y \wedge \hat{y}_{\text{ev}}) \leftrightarrow z$

interactive

$x \rightarrow \text{XPred}(\text{ev.dst}, x, 9:35\text{pm}) \quad x \leftrightarrow y$

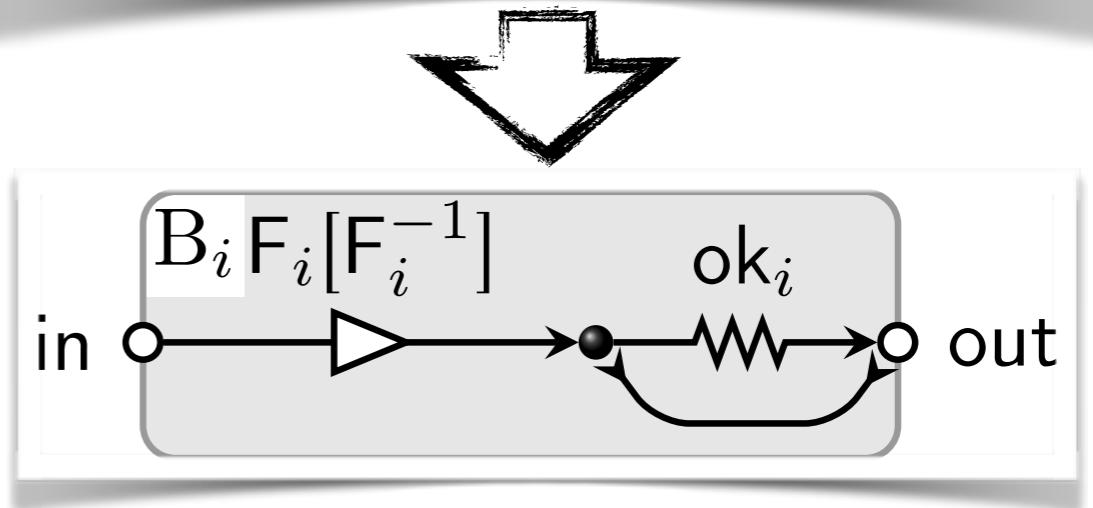
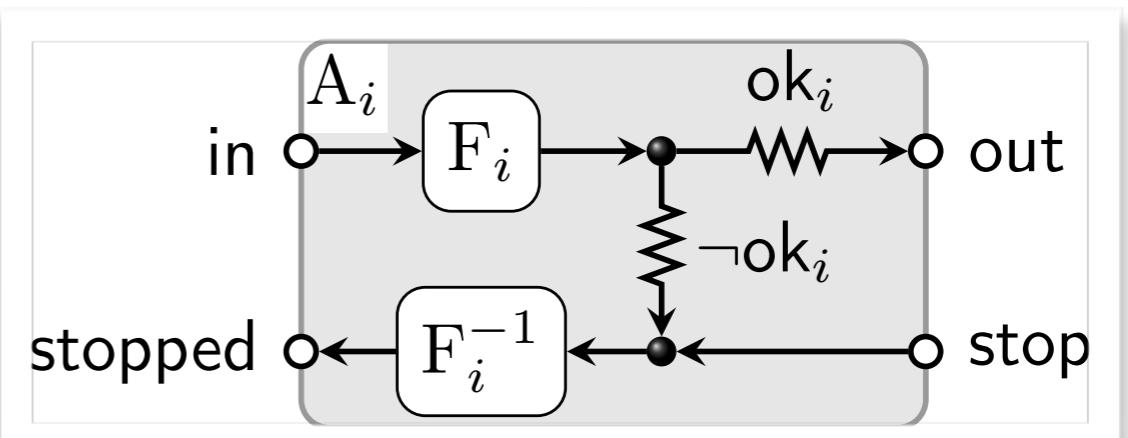
$y \rightarrow \hat{y}_{\text{ev}} := \hat{x}_{\text{ev.dst}} \quad (y \wedge \hat{y}_{\text{ev}}) \leftrightarrow z$

Wrapping up



- Reo and connector colouring

- Interactive constraint solving
- Expose the atomicity of Reo to components



Ongoing experiments

- Avoiding pre-processing (SMT instead of SAT) $\exists a, \hat{a} \cdot \psi$
- Compiling steps:
- Heuristics for identifying potential partial colourings
- Combining local and interactive constraints (probably to a journal)